

Inspira Crea Transforma

The International Spring School on Integrated Operational Problems: What a PhD student can learn

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Doctoral Seminar on Mathematical Engineering - August 3, 2018

Outline

- What is the International Spring School on Integrated Operational Problems – ISSIOP?
- About Troyes in France
- Research and lectures at the ISSIOP
- My schedule at the ISSIOP
- Personal experience and acknowledgements

What is the ISSIOP?

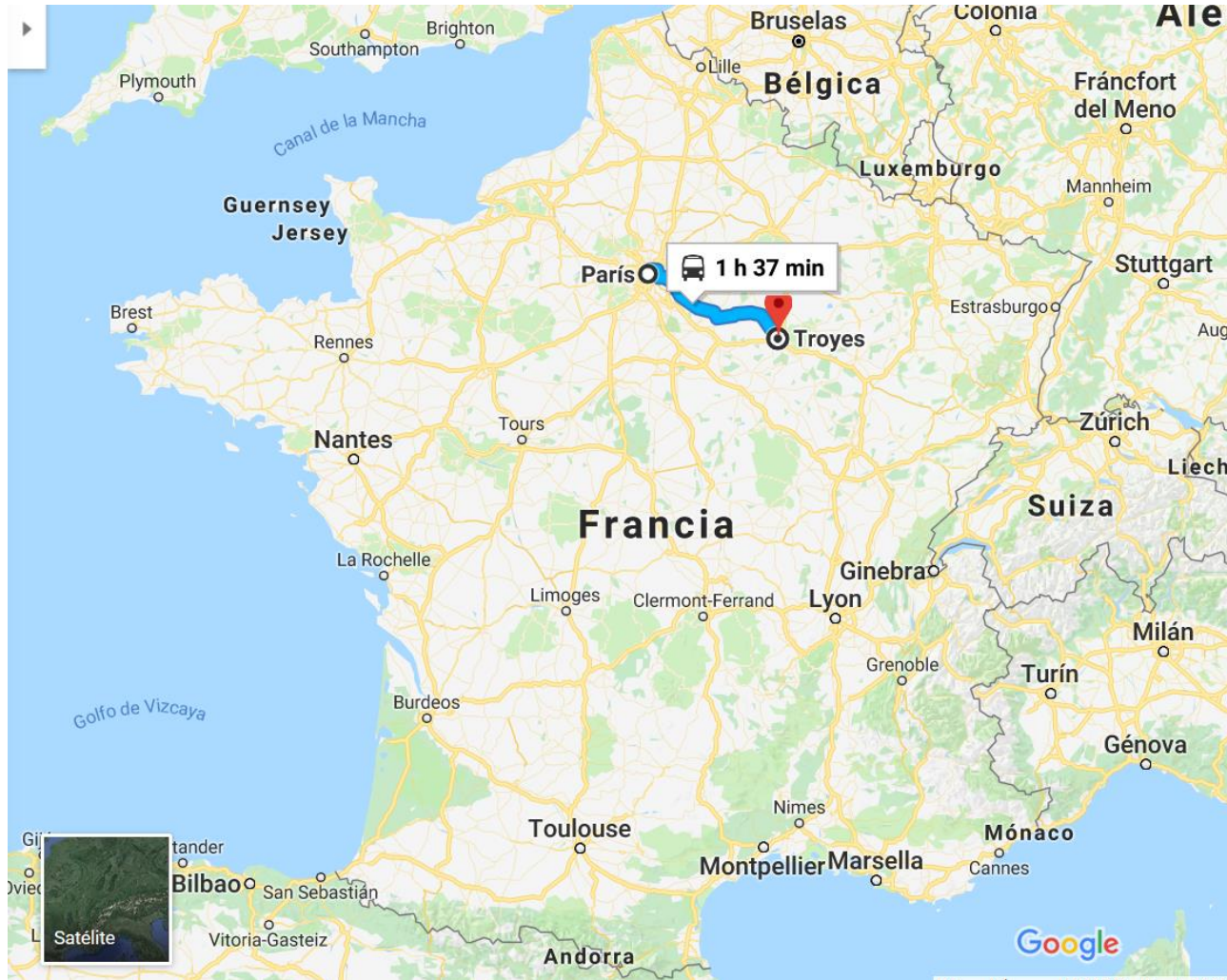
It is the fourth edition of the GdR RO Young Researchers School, organized in collaboration with 4 working groups (GT2L, GT Bermuda, GT Meta, GT OSI)

In 2018 the school takes an international dimension: all the courses are made by scientists from several countries (Spain, Belgium, France).

What is the ISSIOP?



About Troyes in France



About Troyes in France



About Troyes in France



Research and lectures at ISSIOP



Ch. Prins
UTT



J. Billaut
Université F. Rabelais
(Tours)



A. Corberán
Universidad de Valencia



D. Feillet
E. Mines Saint-Etienne



Marc Sevaux
Université de Bretagne-Sud



Kenneth Sörensen
Univ. Antwerpen

Research and lectures at ISSIOP

1. Split algorithms - Ch. Prins
2. Column generation – D. Feillet
3. Integrated scheduling and routing – J. Billaut
4. Introduction to web services – M. Sevaux
5. Optimization of Smart grids – E.G. Talbi
6. Internet of things - L. Merghem-Boulahia
7. The trustfull promises of block chain – Ph. Entzman
8. Constraint programming – E. Hebrard
9. Linear programming for routing - A. Corberán
10. Last advances in metaheuristics – K. Sörensen

My schedule at ISSIOP

Lectures

Split algorithms

Column generation

Integrated scheduling and routing

Constraint programming

Linear programming for routing

Last advances in metaheuristics

Projects

Column generation

Integrated scheduling and routing

Column generation lecture

Vehicle Routing Problem
with Time Windows
(VRPTW)

Column generation
algorithm

Implementation remarks

Ecole Nationale
Supérieure des Mines
SAINT-ETIENNE

www.emse.fr

Spring School - UTT

Column generation and branch-and-price for vehicle
routing problems
Introduction

Dominique Feillet – Mines Saint-Etienne and LIMOS

MINISTÈRE DE L'ÉCONOMIE
DE L'INDUSTRIE ET DE L'EMPLOI

Image taken from: http://fc.isima.fr/~lacomme/Spring_School/conf/slides/

Column generation lecture

subject to

$$\text{minimize } \sum_{1 \leq u \leq U} \sum_{(v_i, v_j) \in A} c_{ij} x_{ij}^u$$

$$\sum_{1 \leq u \leq U} \sum_{\{v_j \in V | (v_i, v_j) \in A\}} x_{ij}^u \geq 1 \quad (v_i \in V \setminus \{v_0\}),$$

$$\sum_{\{v_j \in V | (v_i, v_j) \in A\}} x_{ij}^u - \sum_{\{v_j \in V | (v_j, v_i) \in A\}} x_{ji}^u = 0 \quad (v_i \in V, 1 \leq u \leq U),$$

$$\sum_{\{v_i \in V | (v_0, v_i) \in A\}} x_{0i}^u \leq 1 \quad (1 \leq u \leq U),$$

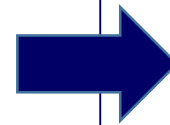
$$\sum_{(v_i, v_j) \in A} d_i x_{ij}^u \leq Q \quad (1 \leq u \leq U),$$

$$s_i^u + st_i + c_{ij} - s_j^u + Mx_{ij}^u \leq M \quad ((v_i, v_j) \in A, v_j \neq v_0, 1 \leq u \leq U),$$

$$s_i^u + st_i + c_{i0} - b_0 + Mx_{i0}^u \leq M \quad ((v_i, v_0) \in A, 1 \leq u \leq U),$$

$$a_i \leq s_i^u \leq b_i \quad (v_i \in V, 1 \leq u \leq U),$$

$$x_{ij}^u \in \{0, 1\} \quad ((v_i, v_j) \in A, 1 \leq u \leq U),$$



subject to

$$\text{minimize } \sum_{r_k \in \Omega} c_k \theta_k$$

$$\sum_{r_k \in \Omega} a_{ik} \theta_k \geq 1 \quad (v_i \in V \setminus \{v_0\}),$$

$$\sum_{r_k \in \Omega} \theta_k \leq U,$$

$$\theta_k \in \mathbb{N} \quad (r_k \in \Omega).$$



The set of all feasible routes

Image taken from: http://fc.isima.fr/~lacomme/Spring_School/conf/slides/

Column generation lecture

Linear relaxation of the extended formulation

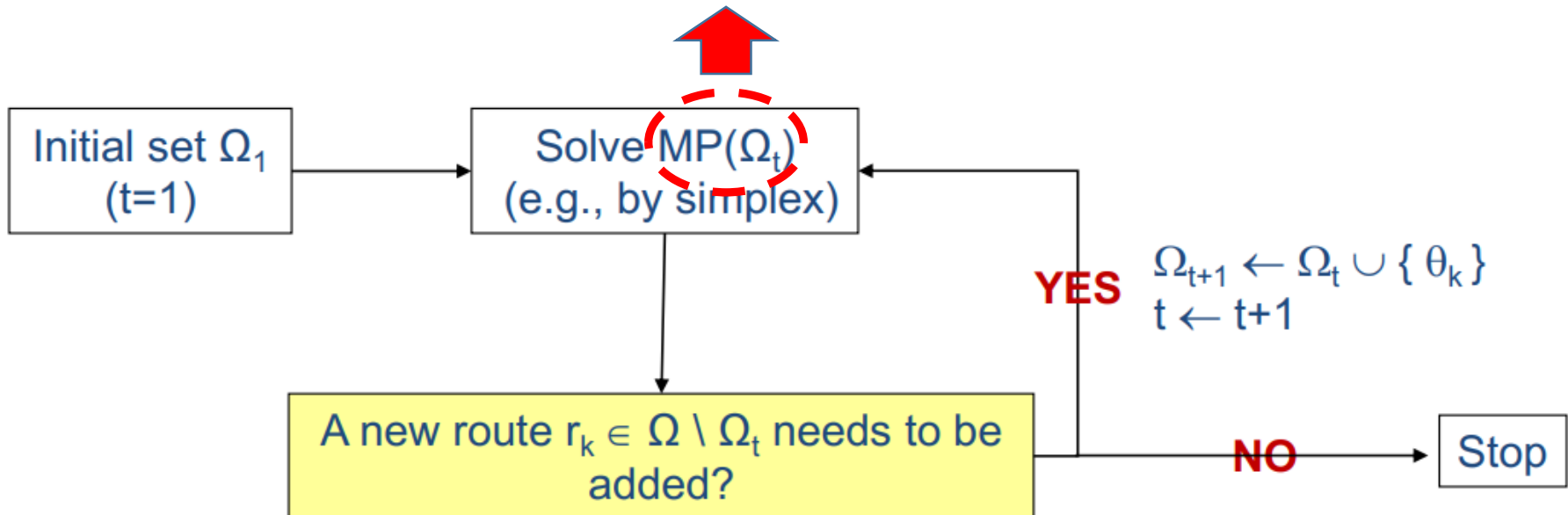


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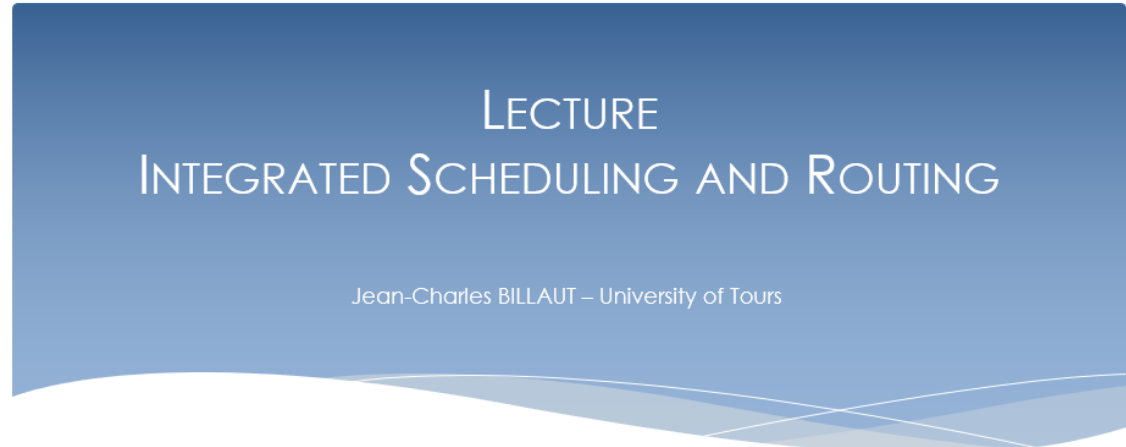
Integrated scheduling and routing lecture

Basics on scheduling and
routing

Integrated problem

GLPK (for a MILP
implementation)

Python (for a metaheuristic
implementation)



International Spring School on Integrated Operational Problems
Troyes
14-16 May 2018

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Integrated scheduling and routing lecture

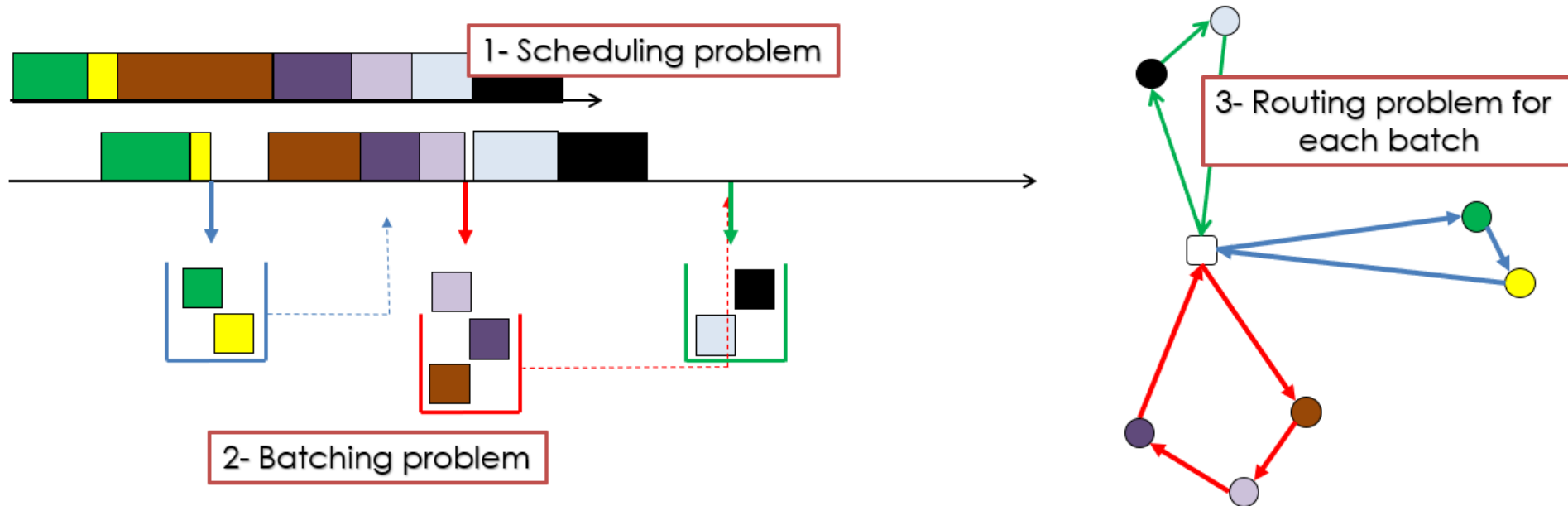


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Integrated scheduling and routing lecture

```
#####  
# scheduling  
#####  
  
s.t. contr_Ck_1 {i in 0..m-1, j1 in 0..n-1, j2 in 0..n-1:j1!=j2}:  
# if j1 precedes j2  
C[i,j2] >= C[i,j1] + p[i,j2] - BigM*(1-y[j1,j2]);  
  
s.t. contr_Ck_2 {j1 in 0..n-1, j2 in 0..n-1:j1!=j2}:  
# j1 before j2 or j2 before j1  
y[j1,j2]+y[j2,j1]=1;  
  
s.t. contr_Ck_2b {i in 1..m-1, j in 0..n-1}:  
# "routing" constraint  
C[i,j] >= C[i-1,j] + p[i,j];  
  
s.t. contr_Ck_2a {j in 0..n-1}:  
# first job  
C[0,j] >= p[0,j];  
  
s.t. def_TjM {j in 0..n-1}:  
# tardiness expression  
Tj[j] >= C[m-1,j] - d[j];  
  
end;
```

MILP implementation for the
scheduling part using GLKP

```
# Data  
#####  
param n;  
param m;  
param p {i in 0..m-1, j in 0..n-1};  
param d {j in 0..n-1};  
param BigM:=sum{i in 0..m-1, j in 0..n-1} p[i,j];  
  
# Variables  
#####  
var y {j1 in 0..n-1, j2 in 0..n-1}, binary; #= 1 if j1 < j2  
var C {i in 0..m-1, j in 0..n-1}, >= 0; # completion time of job j  
var Tj {j in 0..n-1}, >= 0; # tardiness of job j  
  
minimize Obj: sum{j in 0..n-1} Tj[j];
```

Integrated scheduling and routing lecture

```
def InsertTabu(voisinage,indi,indj):
    ElemTabu=[voisinage,indi,indj]
    if len(TabuList) == TabuSize:
        del(TabuList[0])
    TabuList.append(ElemTabu)

def NotTabu(voisinage,indi,indj):
    ElemTabu=[voisinage,indi,indj]
    notTabu = True
    if(ElemTabu in TabuList):
        notTabu = False
    return(notTabu)
```

```
# =====
# Neighborhood
# =====
```

```
def swap_both(i,j,sol):
    # the sequence sol is modified by a swap both in the sequence and in the batches
    #=====
    #print('before swap_seq:',i,j,sol)
    job_i = sol[0][i]
    job_1 = [job_i]
    job_j = sol[0][j]
    job_2 = [job_j]
    sol[0]=sol[0][0:i]+job_2+sol[0][i+1:j]+job_1+sol[0][j+1:njobs]
    u=0
```

```
# =====
# Parameters
# =====

nbjobs=len(pp)
INFINI = 999999999
TabuList=[]
TabuSize = 7
TIME_LIMIT = nbjobs * m / 4
DELTA = nbjobs/2 # is used to limit the swaps
FLAG_SWAP_BOTH = 1
FLAG_SWAP_SEQ = 1
FLAG_SWAP_BATCH = 1
FLAG_2OPT = 0
```

Split algorithms lecture

Route-first cluster-second
methods

Basic splitting procedure

Applications to heuristics
and metaheuristics.



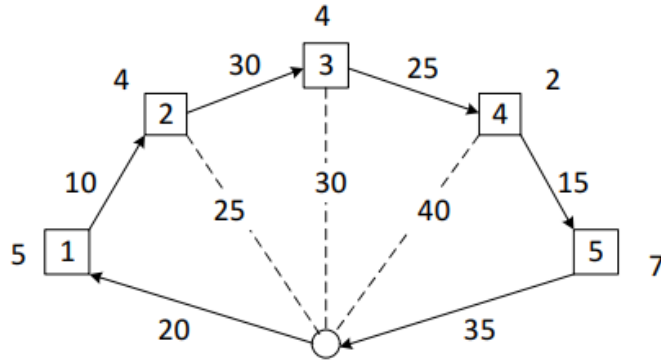
Tour splitting algorithms for vehicle routing problems

Prof. Christian PRINS
christian.prins@utt.fr

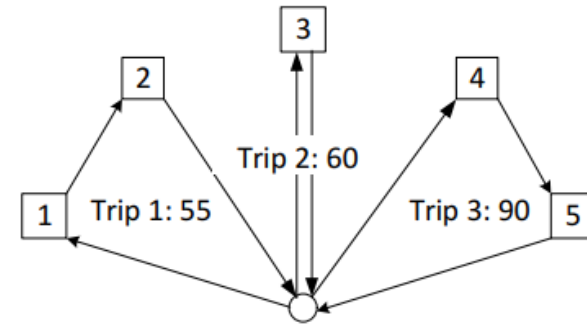
Institute Charles Delaunay (ICD) – UTT
12 rue Marie Curie, CS 42060, 10004 Troyes Cedex, France

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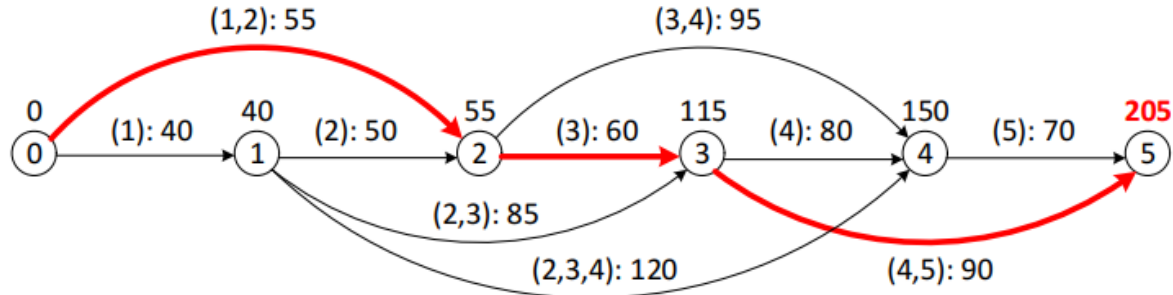
Split algorithms lecture



1. Giant tour $T = (1, 2, 3, 4, 5)$ with demands



3. Optimal splitting, cost 205



2. Auxiliary graph H of possible trips for $Q = 10$ – Shortest path in bold

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Lp-based methods for routing problems

lecture

Polyhedral combinatorics applied to:

- TSP
- OARP
- Close –enough ARP

Lp-based methods for solving routing problems

Ángel Corberán

Universitat de València, Spain

Spring School on Integrated Operational Problems
May 14-16, Troyes, France

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Lp-based methods for routing problems

lecture

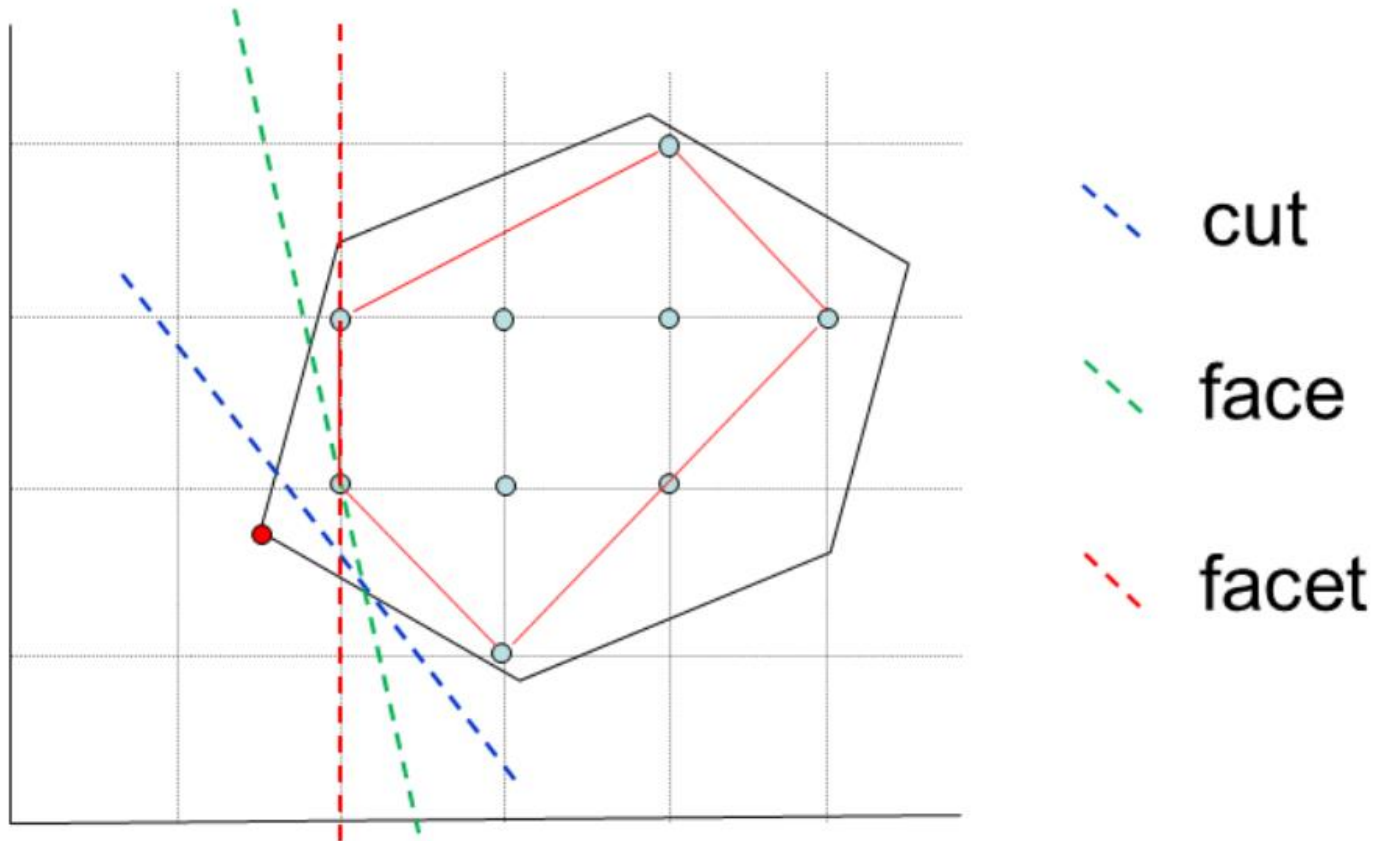


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Last advances in metaheuristics lecture

Complementary local
search operators

Properties of good solutions

Solution metrics

Designing a heuristic the modern way

Or: how to solve very large vehicle routing problems

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International Spring School on Integrated Operational Problems -
Troyes - 14-16 may 2018



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Last advances in metaheuristics lecture

State of the art



- Use as many local search (constructive) operators as possible
- Either VNS or LNS
- Fit in a metaheuristic framework
 - This is your Unique Selling Point
 - But it really does not matter all that much
- Beware of “Frankenstein” algorithms

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Last advances in metaheuristics lecture

Comparison to other algorithms

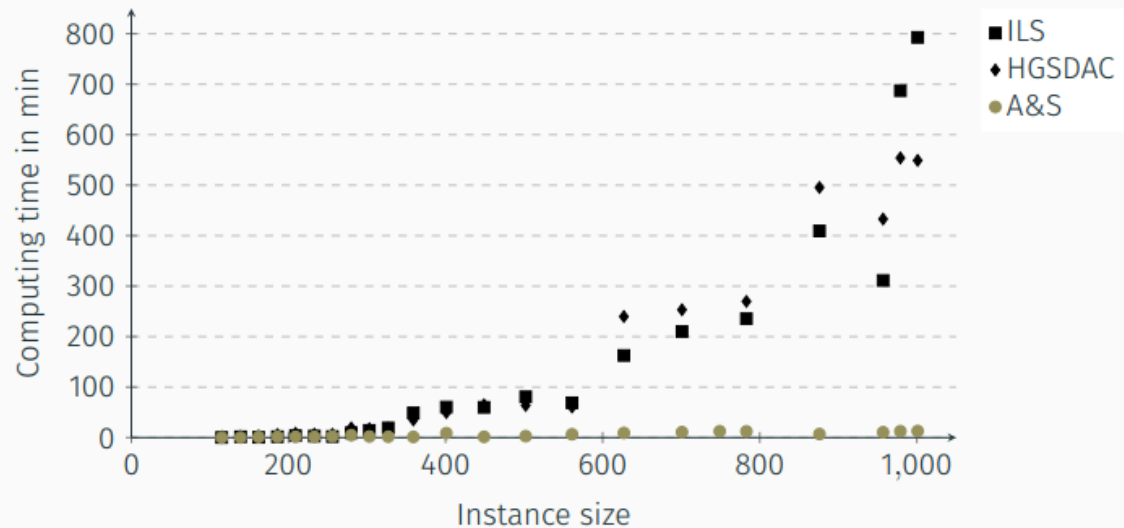


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Personal experience and acknowledgements



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